

# Computer Networks

## X\_400487

Lecture 10

Course Recap



Lecturer: Jesse Donkervliet



WAN Types

Dynamic IP  
Static IP

PPPoE  
PPTP  
L2TP

NAT Forwarding

Port Forwarding  
Port Triggering  
DMZ  
UPnP  
Virtual Server

IGMP Proxy

# 1. After this course, you understand router specifications

## WIRELESS

Standards

IEEE 802.11ax 6 GHz,  
IEEE 802.11ax/ac/n/a 5 GHz,  
IEEE 802.11ax/n/b/g 2.4 GHz

DHCP

Address Reservation  
DHCP Client List  
Server

WiFi Speeds

**AXE5400**  
6 GHz: 2402 Mbps (802.11ax)  
5 GHz: 2402 Mbps (802.11ax)  
2.4 GHz: 574 Mbps (802.11ax)

DDNS

TP-Link  
NO-IP  
DynDNS

Working Modes

Router Mode  
Access Point Mode

WiFi Capacity

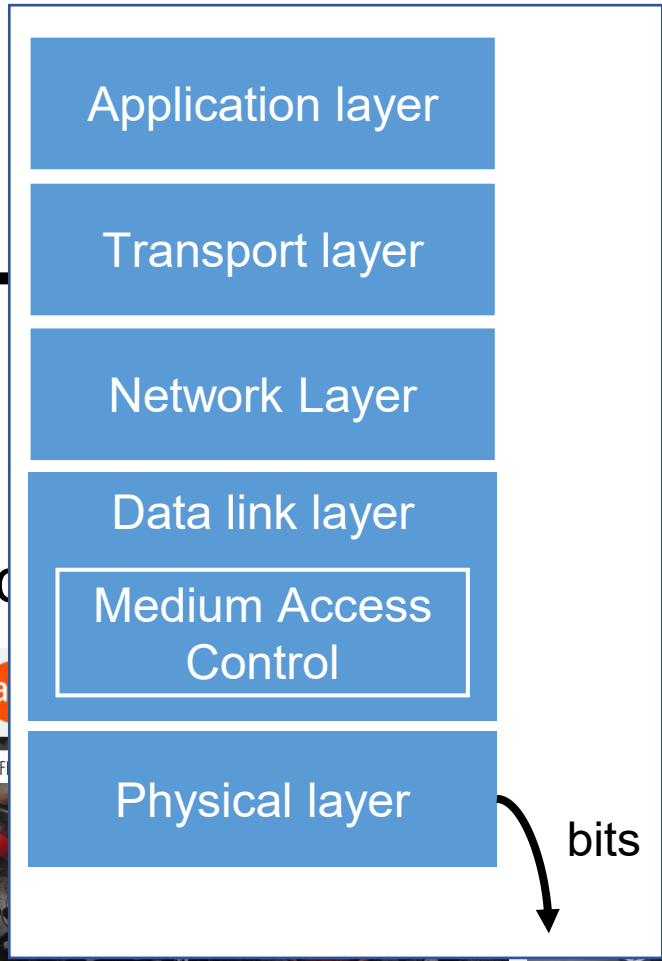
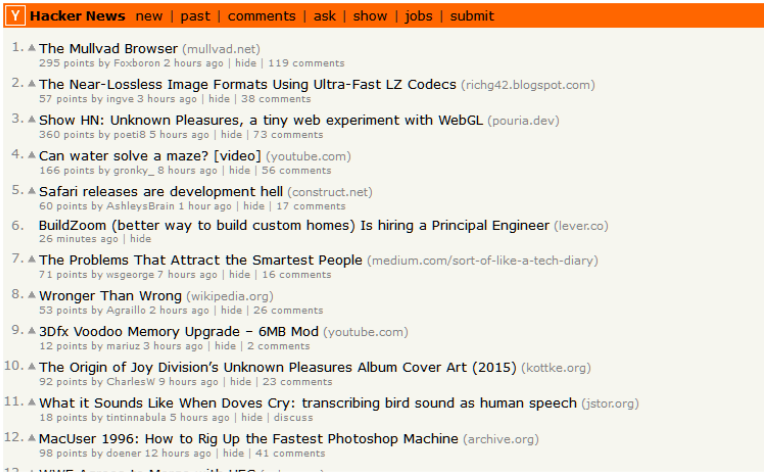
OFDMA

Simultaneously communicates with multiple Wi-Fi 6 clients

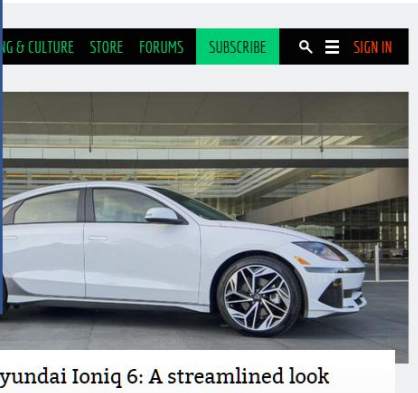
Airtime Fairness

Improves network efficiency by limiting excessive occupation

# (a) Type address in browser



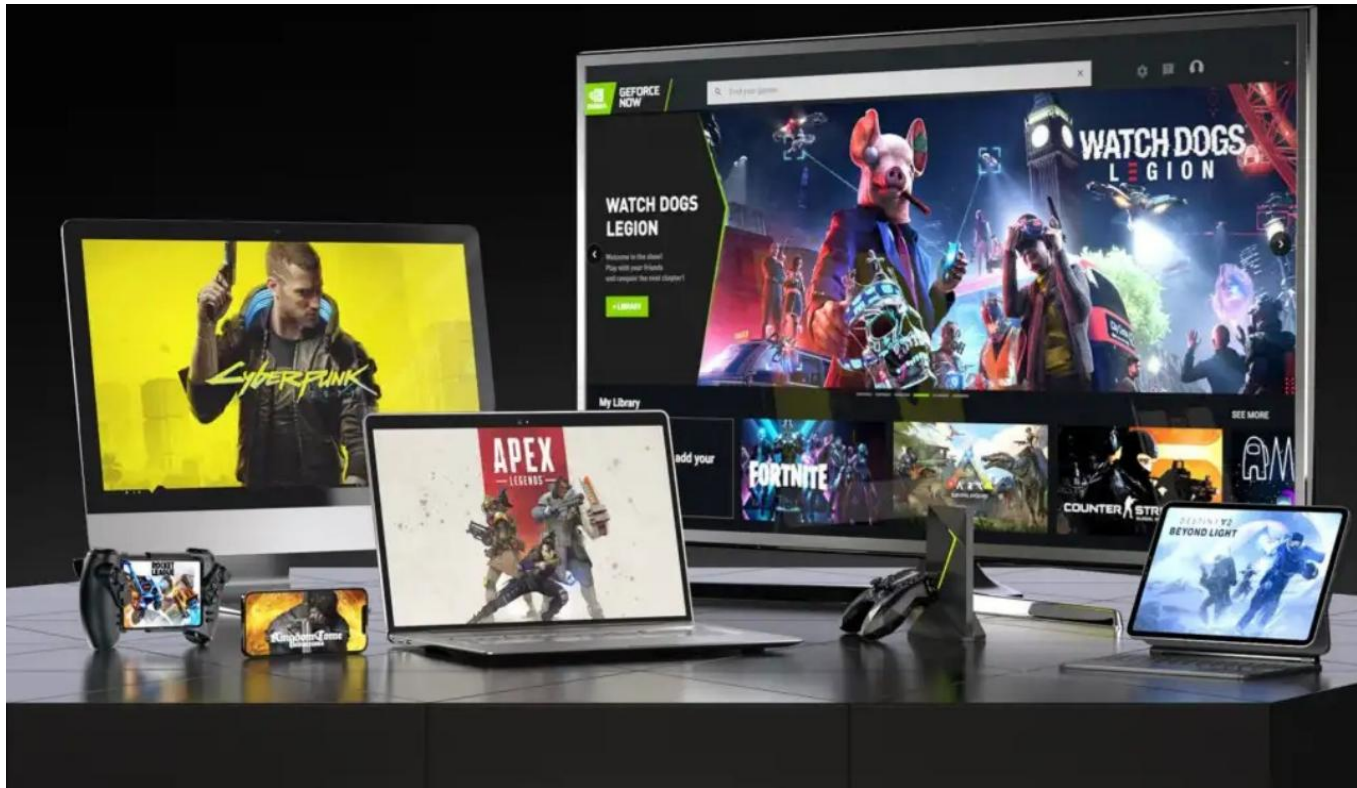
# (b) Browser does magic



Stable Diffusion copyright lawsuits could be a legal

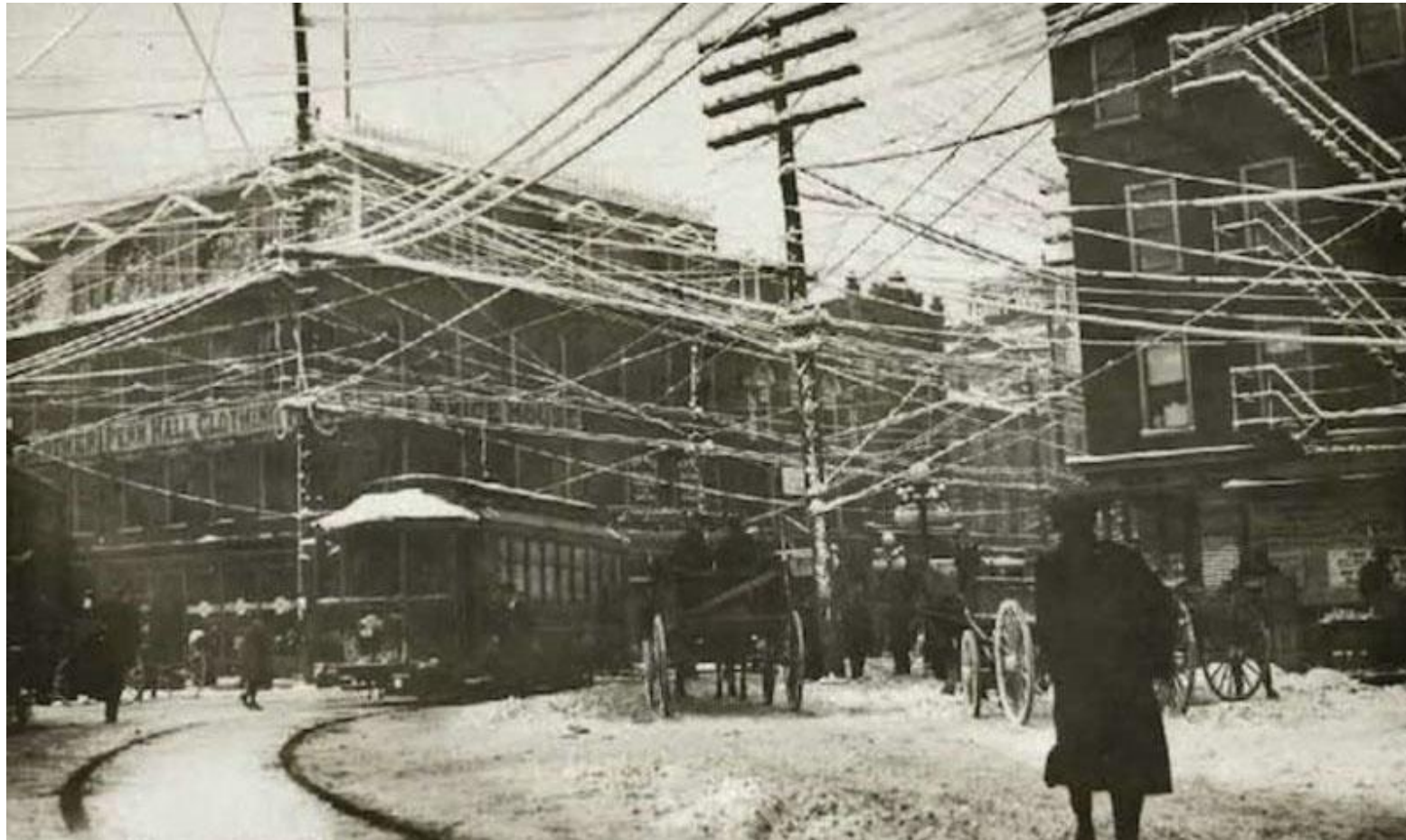
The 2023 Hyundai Ioniq 6: A streamlined look

2. After this course, you understand browser magic



3. After this course, you understand how networks enable new applications

# Early telephone system



# THE INTERNET IN 2023 EVERY MINUTE



Created by: eDiscovery Today & LTMG

## Google Services Go Down in Some Parts of U.S.

People experienced outages of services like Gmail, YouTube and Google Meet.

## Facebook's outage likely cost the company over \$60 million

Configuration change cascaded down the data centers, bringing systems to a halt.

AX SHARMA - 10/5/2021, 2:33 PM

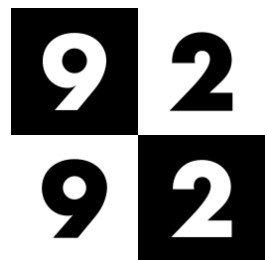
TECH AMAZON

## An Amazon server outage caused problems for Alexa, Ring, Disney Plus, and deliveries

Amazon says "many services have already recovered"

By Richard Lawler | @rjcc | Updated Dec 7, 2021, 7:25pm EST

- Further reading on Facebook outage:
- <https://blog.cloudflare.com/october-2021-facebook-outage/>
  - <https://engineering.fb.com/2021/10/05/networking-traffic/outage-details/>



Rabobank



canvas



Blackboard

# TODAY'S AGENDA

1. Intro
2. Content recap & exercises
3. Game demos
4. Wrap-up

# TODAY'S AGENDA

1. **Intro**

2. Content recap & exercises

a. **Physical Layer**

b. Link Layer

c. MAC Layer

d. Network Layer

e. Transport Layer

f. Application Layer

3. **Game demos**

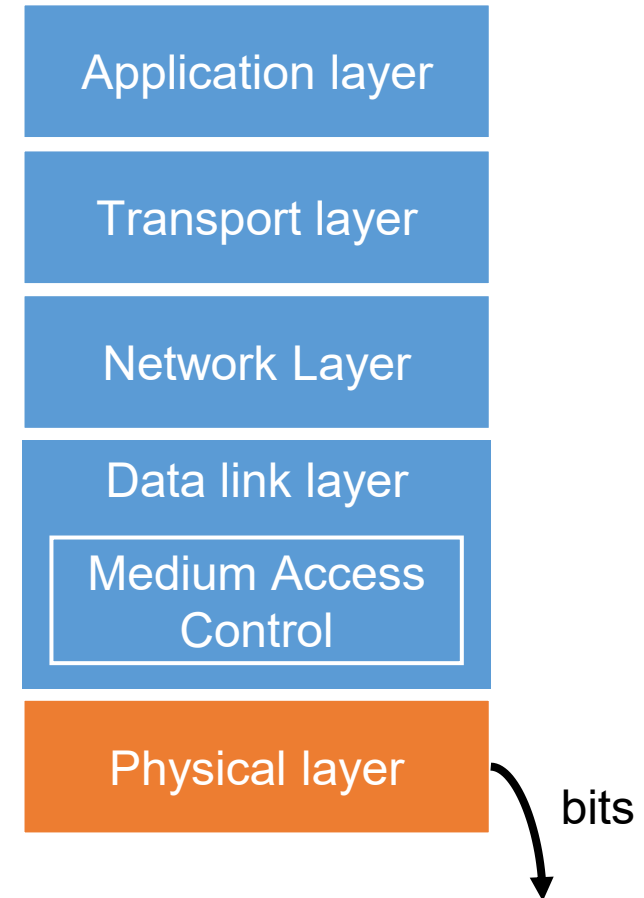
4. **Wrap-up**

# Responsibilities of the Physical Layer

1. Translating between bits/symbols and signals
2. (Static) channel allocation

Important properties:

1. Bit rate
2. Delay
3. Storage capacity
4. Error rate



# Nyquist & Shannon's theorems

## Nyquist's theorem

Computing the maximum data rate for a *noiseless* channel.

$$R = 2B \times \log_2(V)$$

R - maximum data rate

B - bandwidth

V - number of signal levels

## Shannon's theorem

In practice, *noise* reduces the maximum data rate.

$$R = B \times \log_2(1 + S/N)$$

Signal to noise ratio (SNR) is expressed in decibel.

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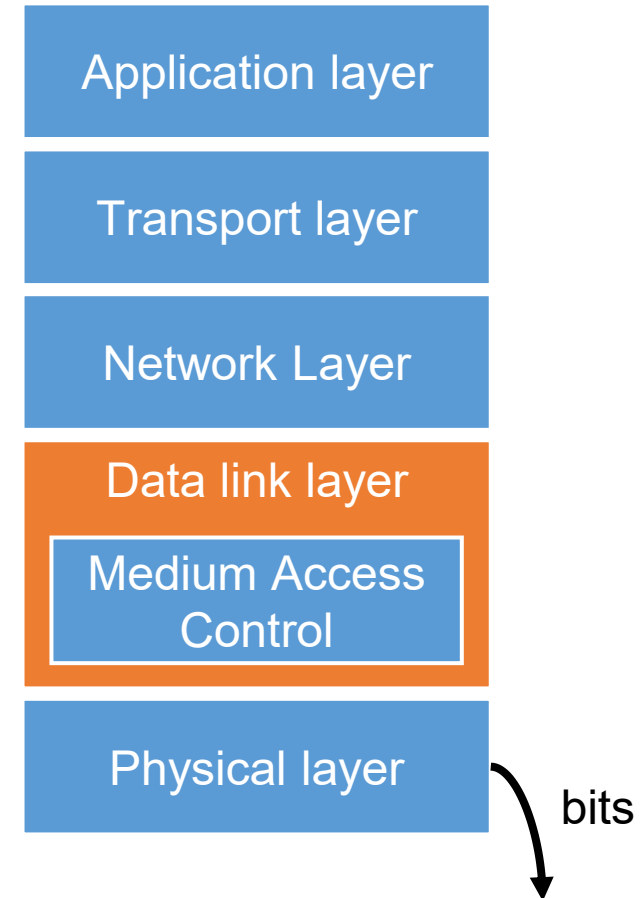
4. **Wrap-up**

# Responsibilities of the Data Link Layer

1. Framing
2. Flow control
3. Error control

Important properties:

1. Bandwidth efficiency
  - Sliding window size
  - Code rate
2. Number/types of errors reliably detected



# Cyclic Redundancy Check Example

Sender adds CRC

$$1 \times x^4 + 0 \times x^3 + 0 \times x^2 + 1 \times x^1 + 1 \times x^0$$

message: 110101010000

generator: 10011

10011010000

10011

10000

10011

0011

Message: 11010101, CRC: 0011,

Codeword: 110101010011

$$x^4 + x + 1$$

Modulo 2 arithmetic.  
No carries/borrows

Q: Consequences for  
implementation?

$$\begin{array}{r} 110101010011 \\ \hline 10011 \\ \hline = 0 \end{array}$$

# Hamming codes - an example

Use bit-locations that are a power of 2 as check bits. Use the remaining positions for the message.

message:        1 101 0101

codeword: \_ \_ 1 \_ 101 \_ 0101

positions:    1 2 3 4 5 6 7 8 9...

1. Expand all bit locations into powers of two.
2. Decide the value of each check bit in position  $2^i$  by calculating the parity function over all bits that have  $2^i$  in their expansion.

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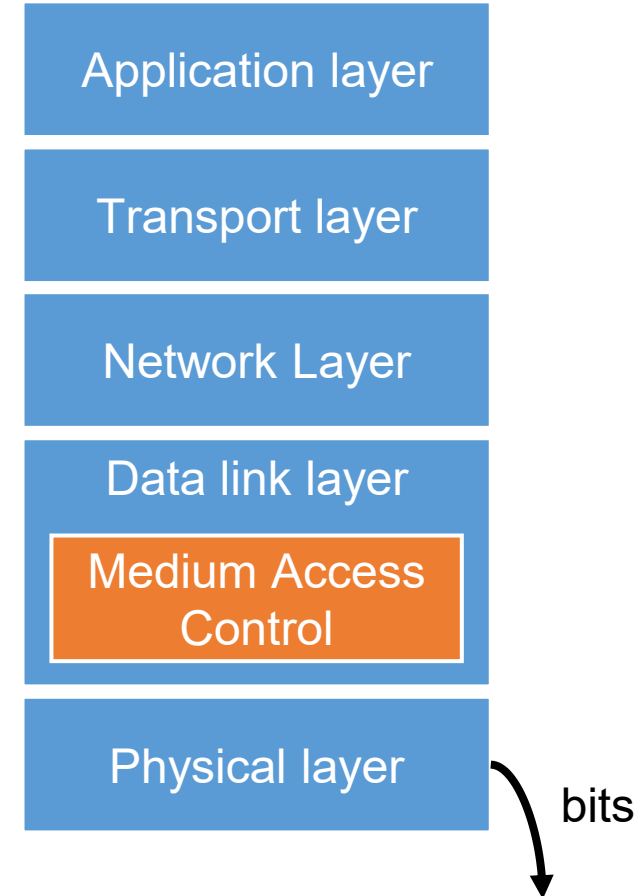
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# Responsibilities of the MAC Layer

1. (Dynamic) channel allocation
2. Collision detection/avoidance
3. Quality of Service

Important properties:

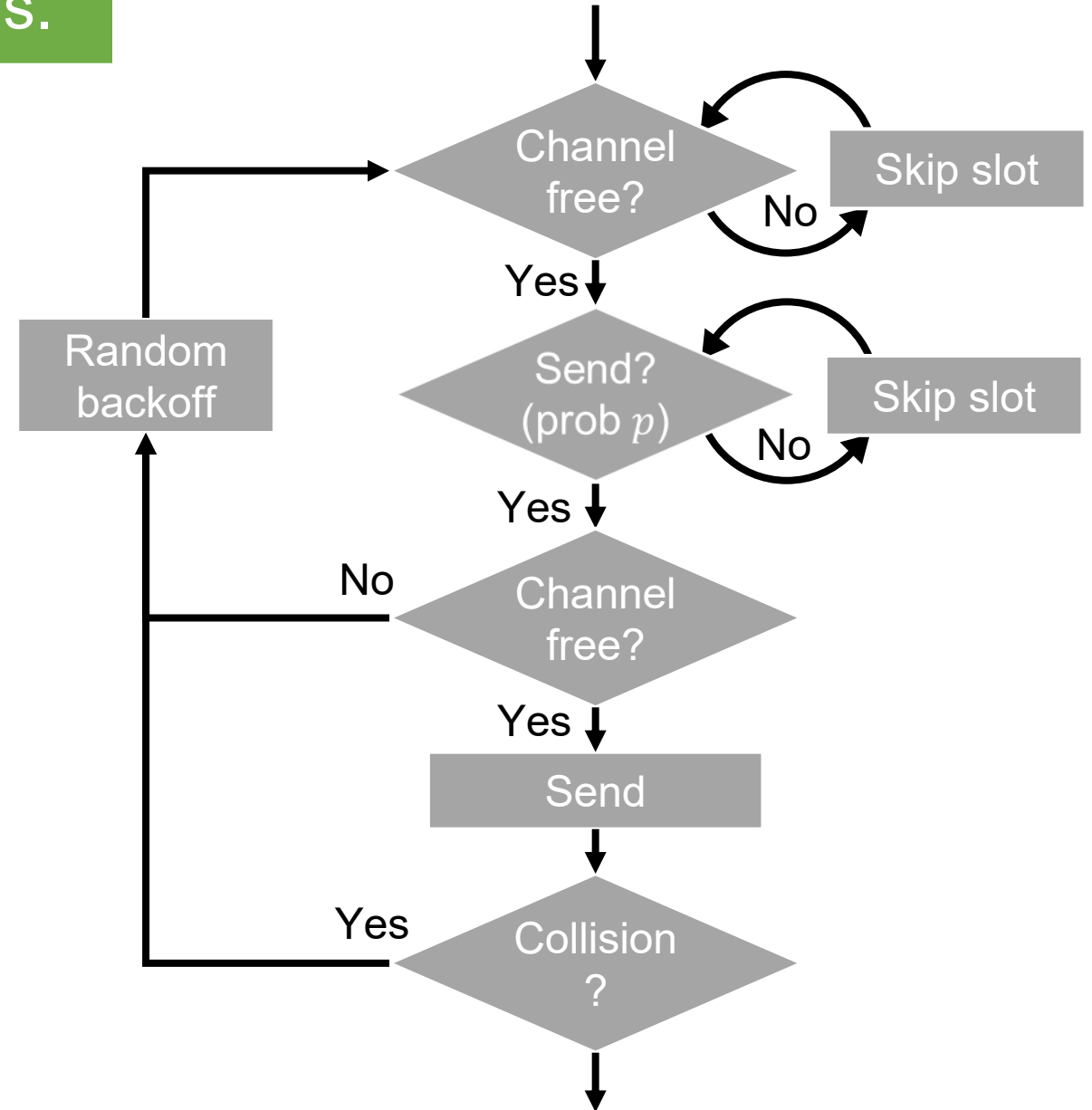
1. Bandwidth efficiency as function of #stations

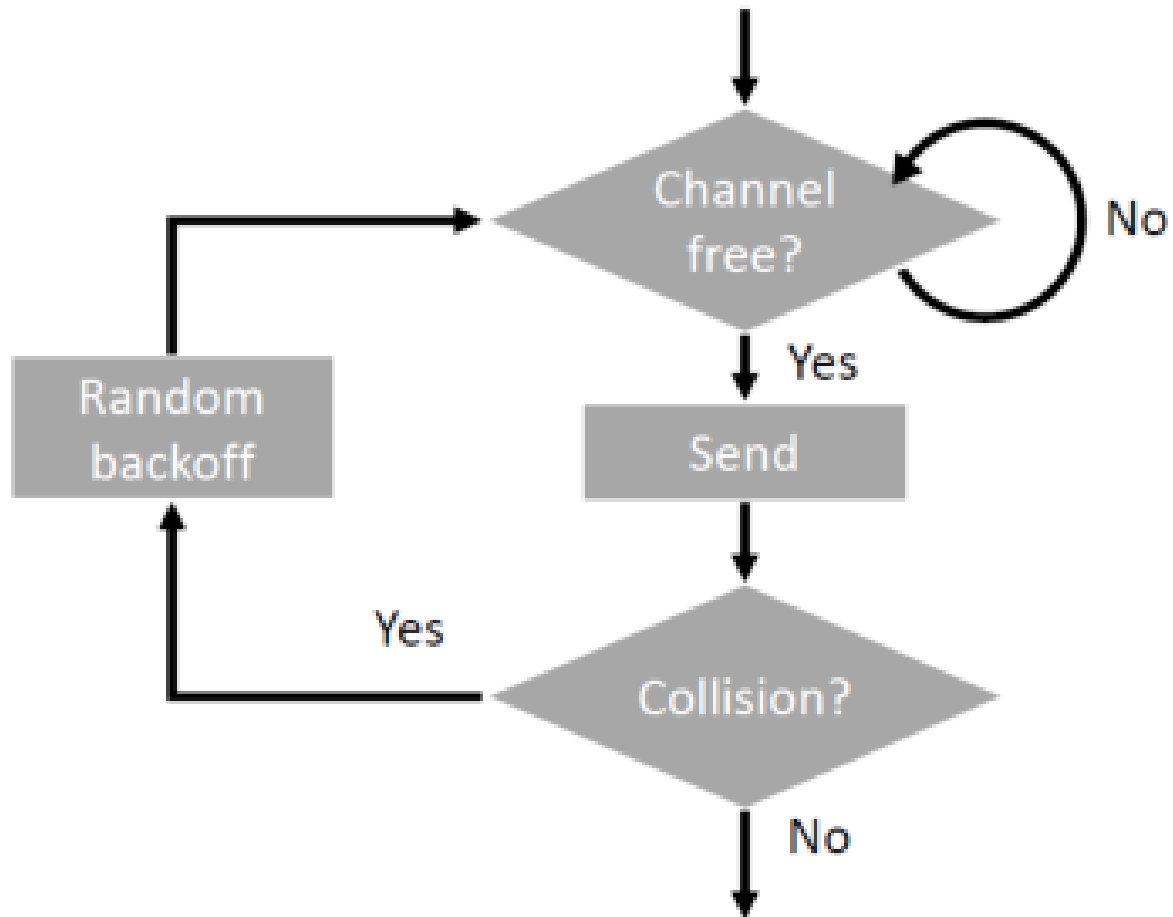


Applies to *slotted* channels.

# $p$ -persistent Carrier Sense Multiple Access (CSMA)

Keeps waiting. Sends frame with probability  $p$ .

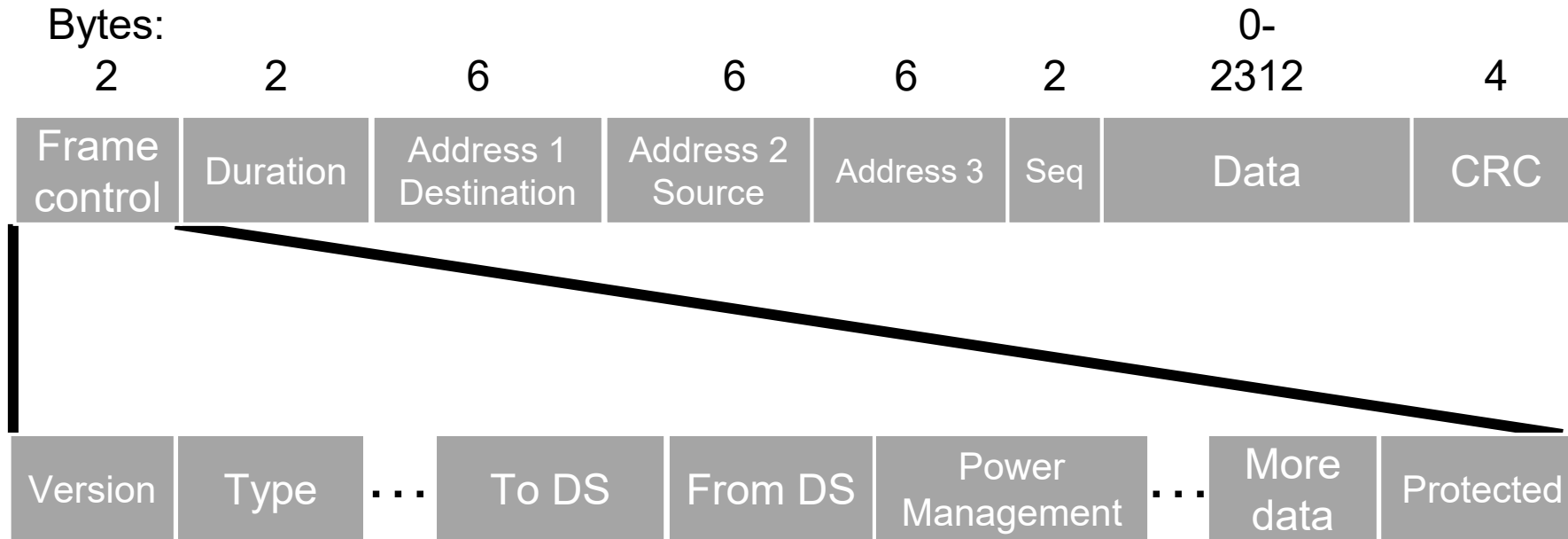




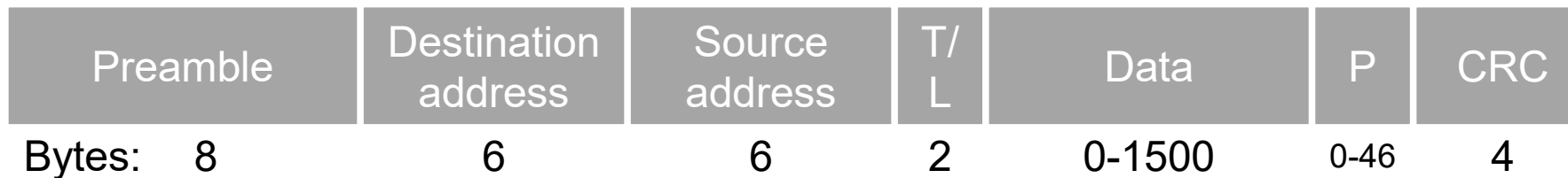
Q: Which protocol is this?

A: 1-persistent CSMA

# 802.11 frame

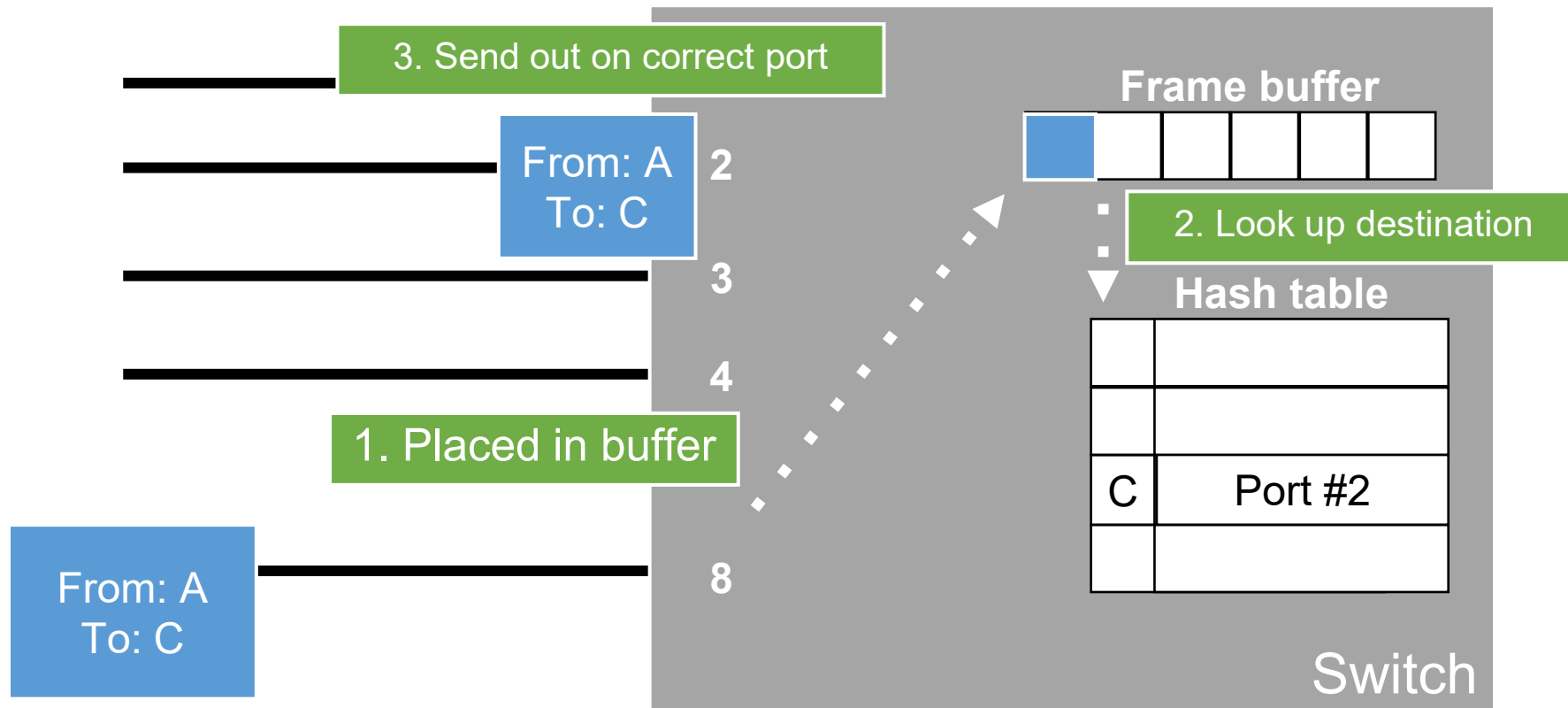


# Ethernet frame



# Ethernet switch

Q: Advantages of switches?



# The Network Layer

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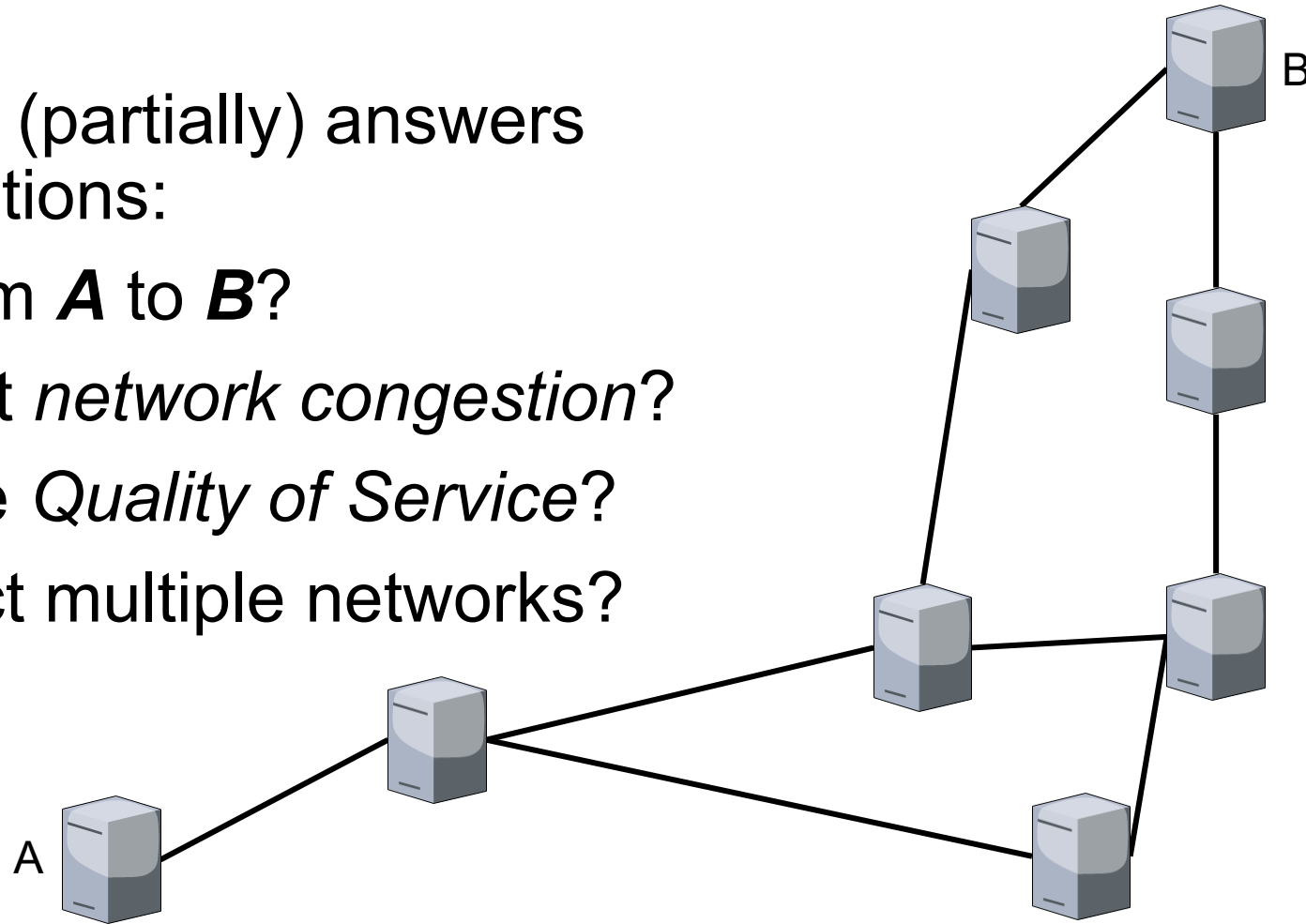
4. **Wrap-up**

# The Network Layer

Lowest layer concerned with end-to-end delivery.

The network layer (partially) answers the following questions:

1. How to get from **A** to **B**?
2. How to prevent *network congestion*?
3. How to provide *Quality of Service*?
4. How to connect multiple networks?

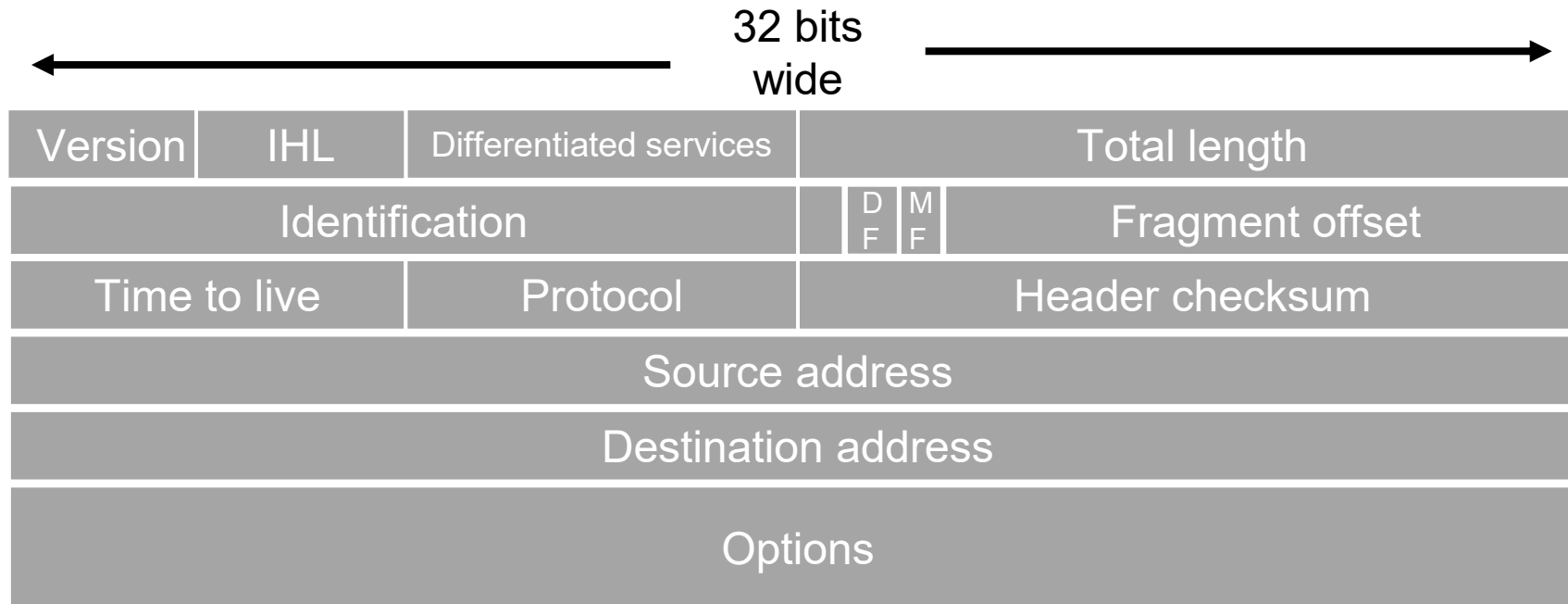


# IP version 4

Q: What is the maximum size of an IP packet?

Q: If IHL field has value 6, how long is the IP packet header?

Frame header: 20-60 bytes (Options 0-40 bytes)



Check the book for the detailed view!

# Routing tables

You want to know for every address, on which link to forward the packet. For this we use a routing table.

Routing table for C.

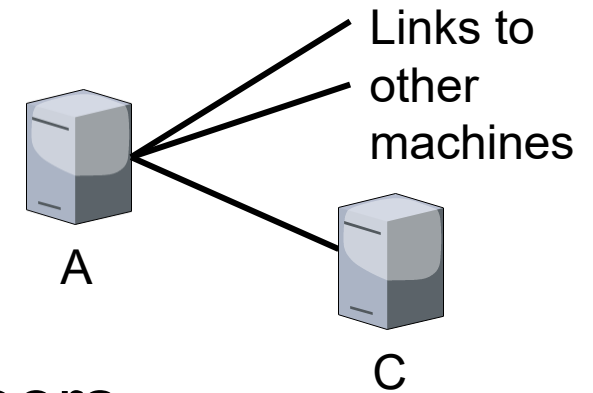
We are directly connected to A

We are station C

To	Distance	Line
A	7	A
B	59	A
C	0	-
D	75	E
E	1	E
F	103	F

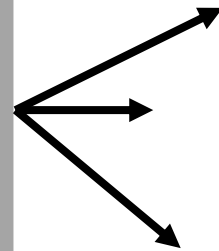
If packet should go to D,  
minimum distance is 75.  
Forward to E.

# Distance vector routing



1. Send your *distance vector* to your neighbors.
2. You use incoming *distance vectors* from your neighbor to construct a *routing table*.

*Distance vector A*  
A, 0  
B, 1  
C, 7  
D, 152  
E, 8  
F, 110



Routing table for C.

To	Distance	Line
A	7	A
B	8	A
C	0	
D	75	
E	1	
F	103	

Q: What can go wrong when using this algorithm?

Routing algorithms can calculate routes to prefixes, instead of to every individual address

# Internet Protocol Prefixes and Subnets

Vrije Universiteit given a *prefix*. E.g., all IP addresses that match **37.60.x.y**.

16 bits used by network

Address starts with 37.60?  
If yes, route to VU.

Example address: **37.60.194.64**.

**00100101.00111100**.11000010.01000000

Network

Host

Prefix: 37.60.0.0/16

Subnet mask:

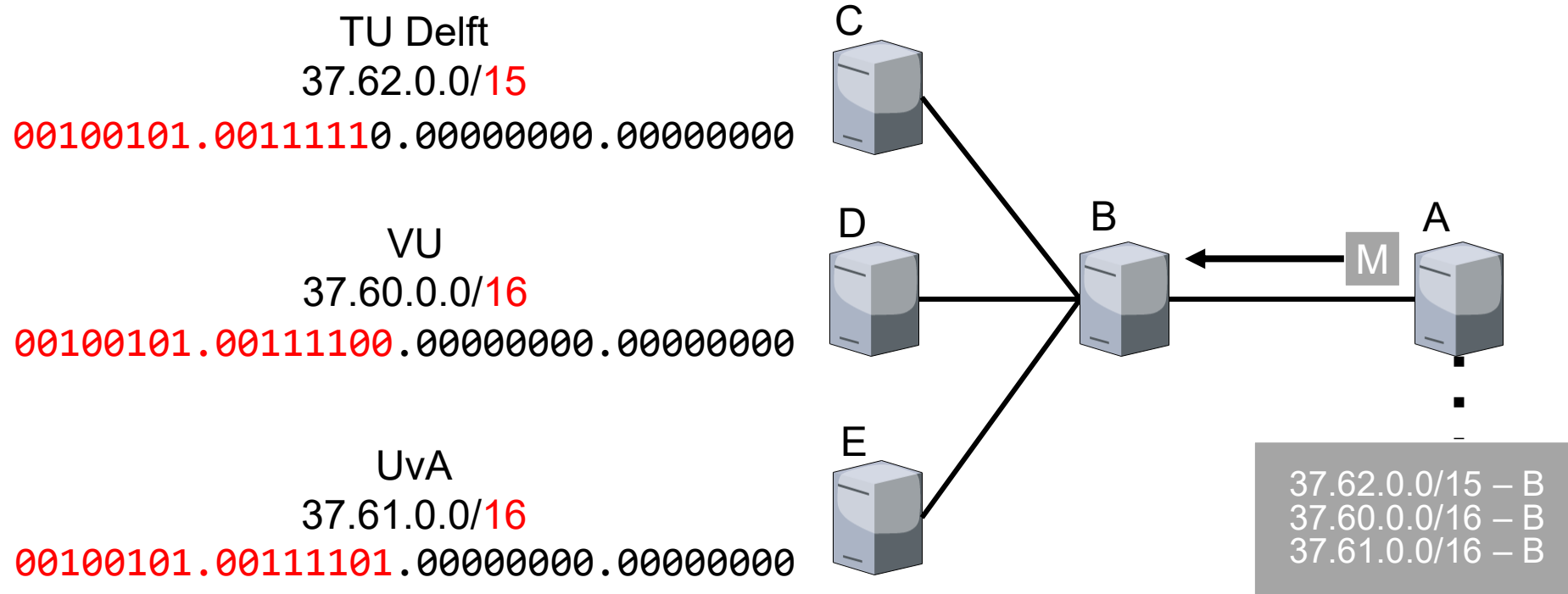
11111111.11111111.00000000.00000000

Prefixes handed out by  
single organization: ICANN

Organizations can  
further subdivide their  
prefix to create *subnets*

# Internet Protocol - CIDR

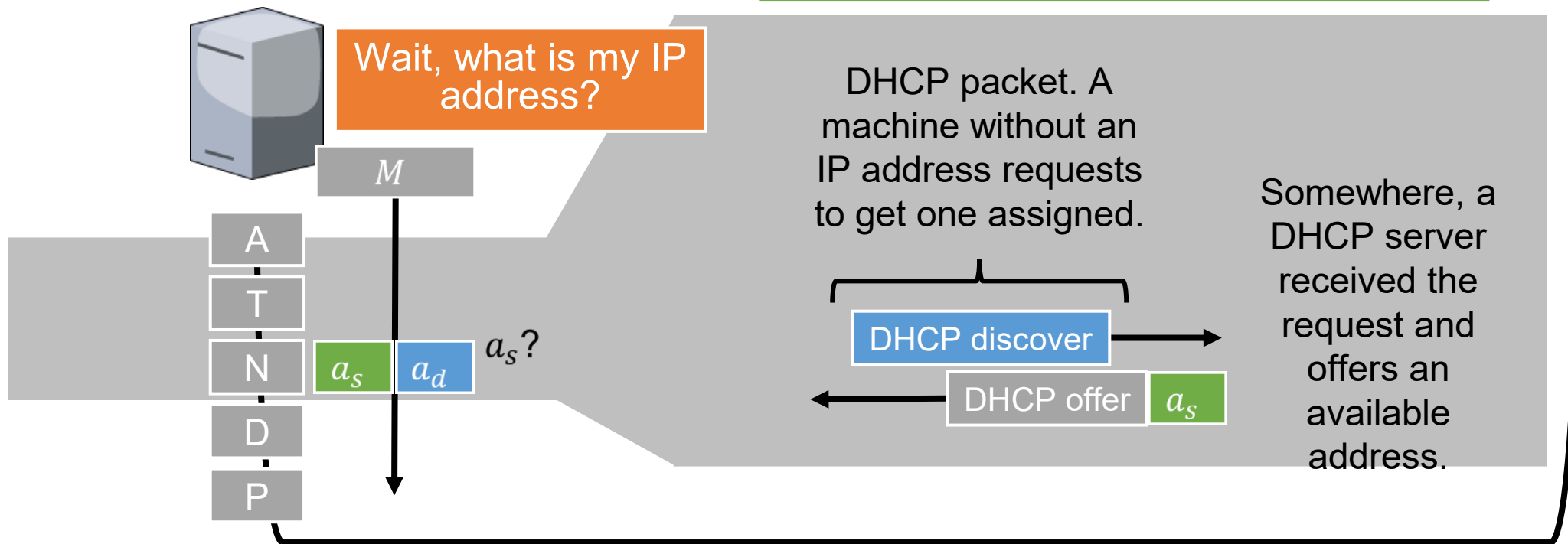
## Classless InterDomain Routing



# Dynamic Host Configuration Protocol (DHCP)

MAC addresses are built into NICs. But network addresses are not.

Used to configure other settings such as: **DNS name servers**, addresses of default gateway, time servers, etc.



Q: How to send DHCP offer back to machine without an address?

# Network Address Translation (NAT)

$a_s$  = source address

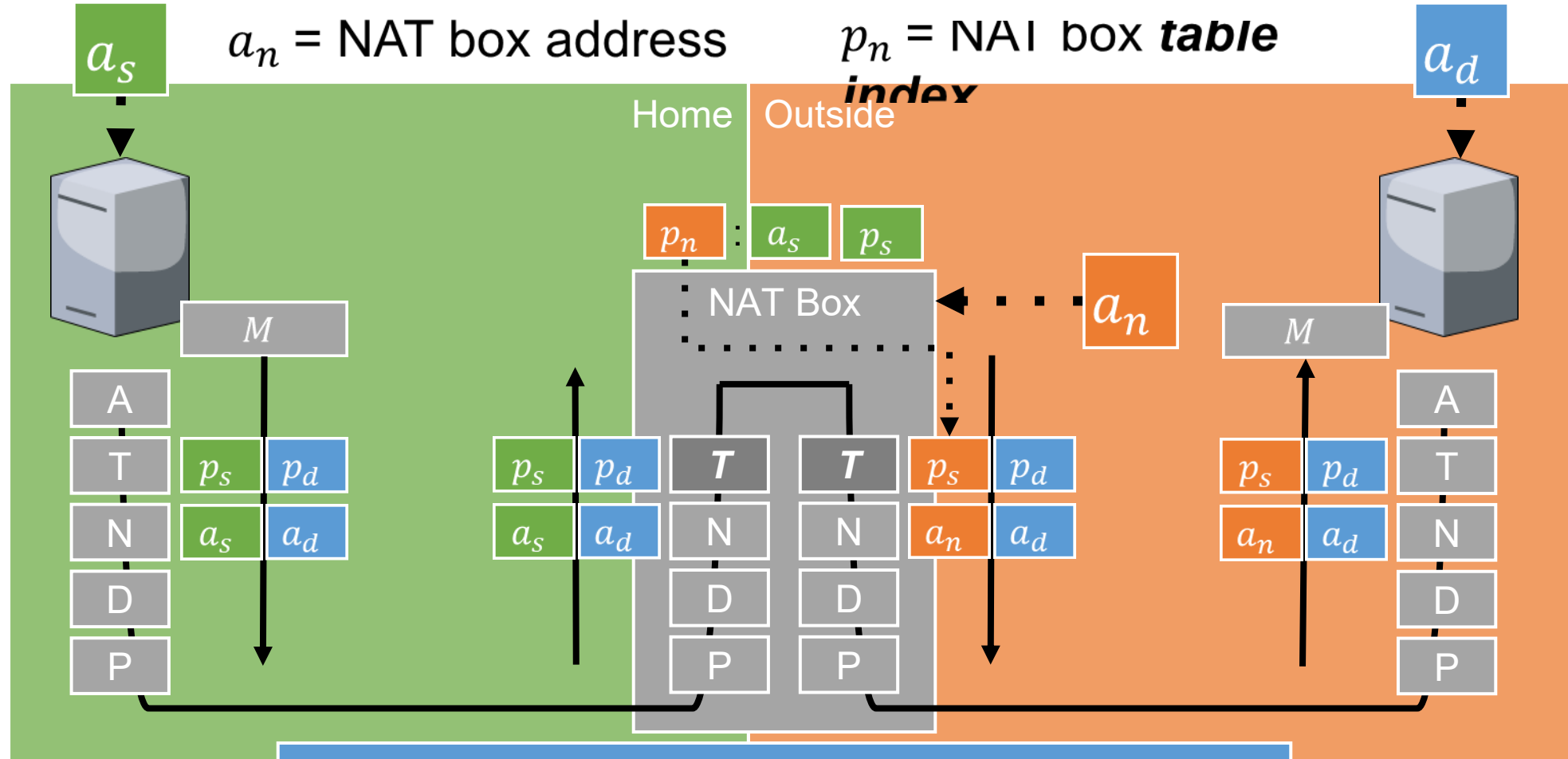
$p_s$  = source port

$a_d$  = destination address

$p_d$  = destination port

$a_n$  = NAT box address

$p_n$  = NAT box **table**



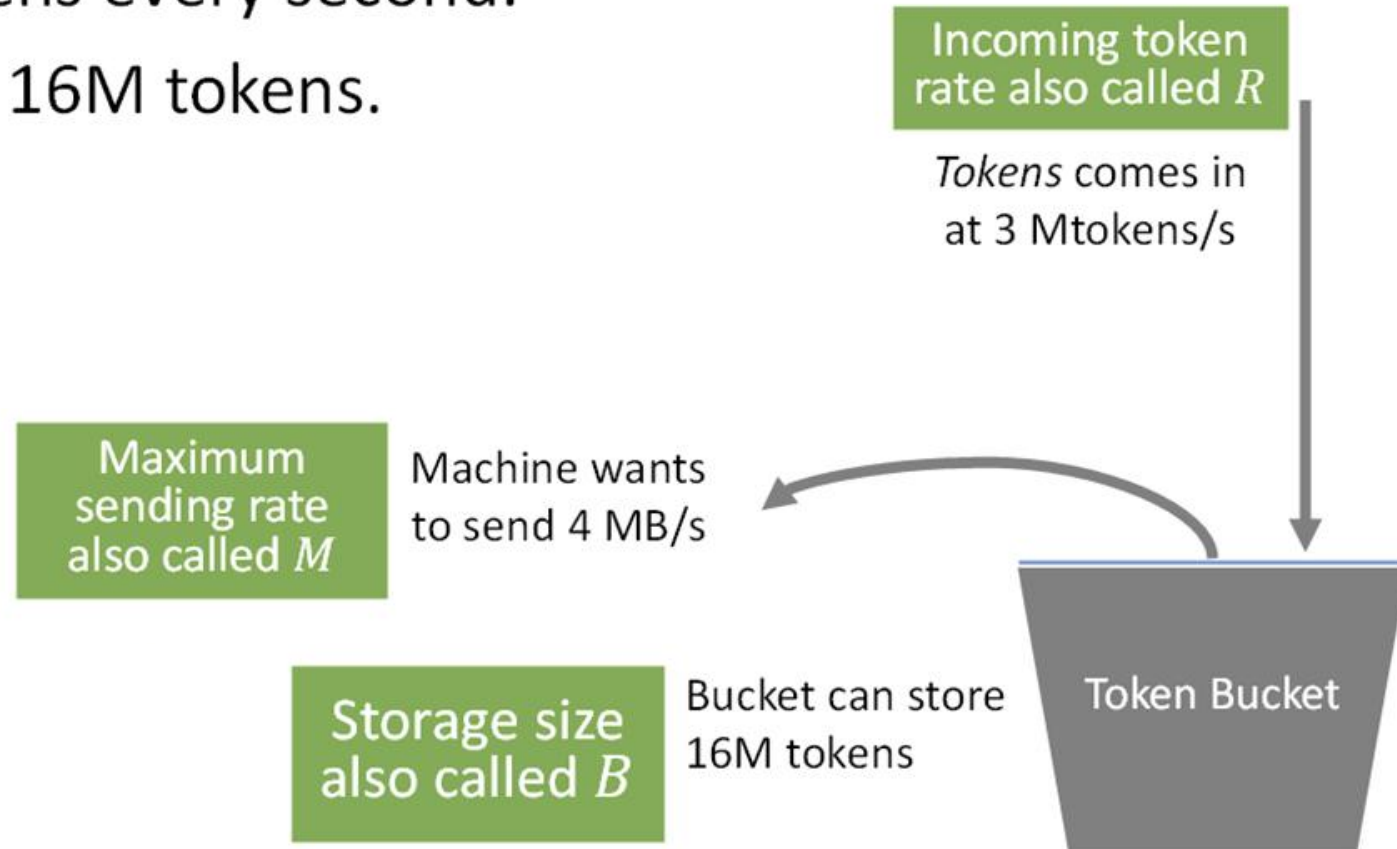
Q: Objections to this approach?

# Traffic shaping

## Token bucket example

Maximum *burst duration* is  $\frac{B}{M-R}$  seconds

Bucket loses 1Mtokens every second.  
Full bucket contains 16M tokens.



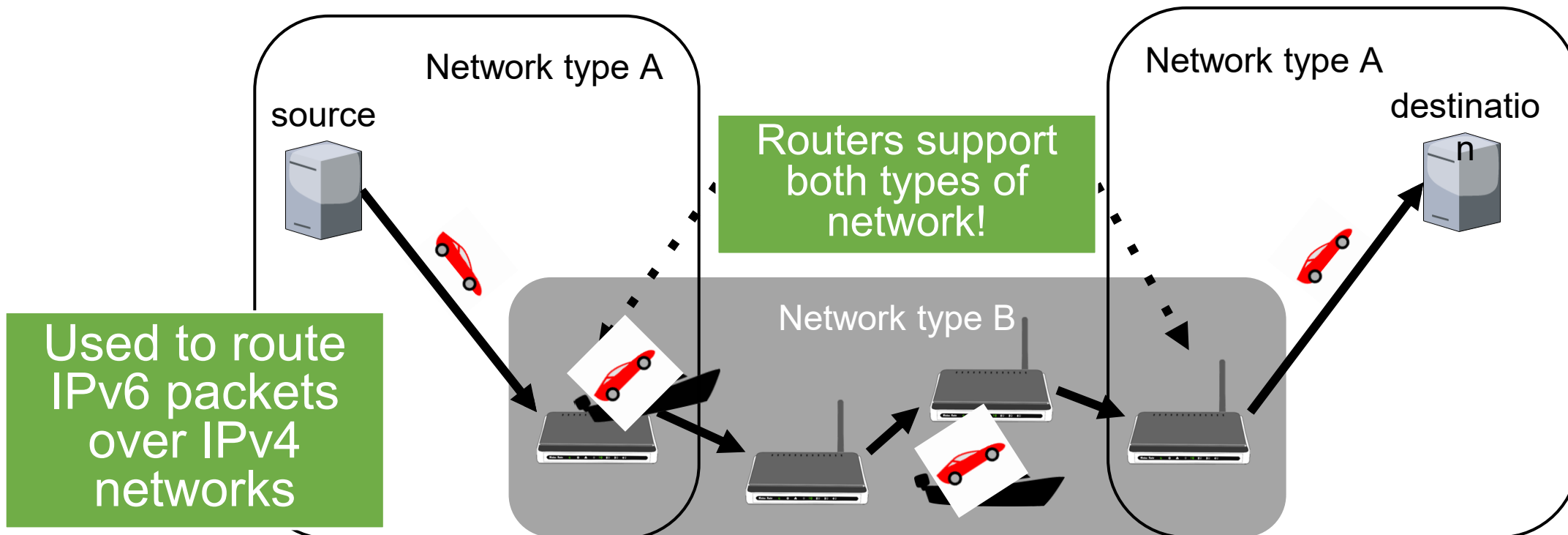
# Tunneling

Q: How to route IPv6 traffic over IPv4 network?

A: [IPv6 hdr][IPv4 hdr][data]

B: [IPv4 hdr][IPv6 hdr][data]

If the **source** and **destination** networks use the same protocols, we can use **tunneling**.



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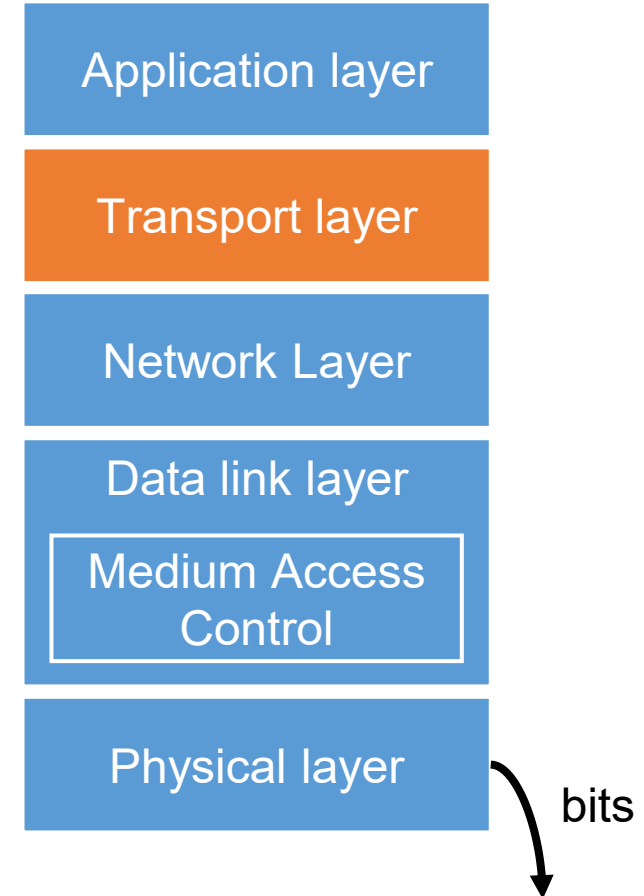
f. Application Layer

3. Game demos

4. **Wrap-up**

# Responsibilities of the Transport Layer

1. Fair and dynamic bandwidth allocation
2. Prevent network congestion
3. Provide (reliable) network service to applications
4. Support Quality of Service



# The transport layer

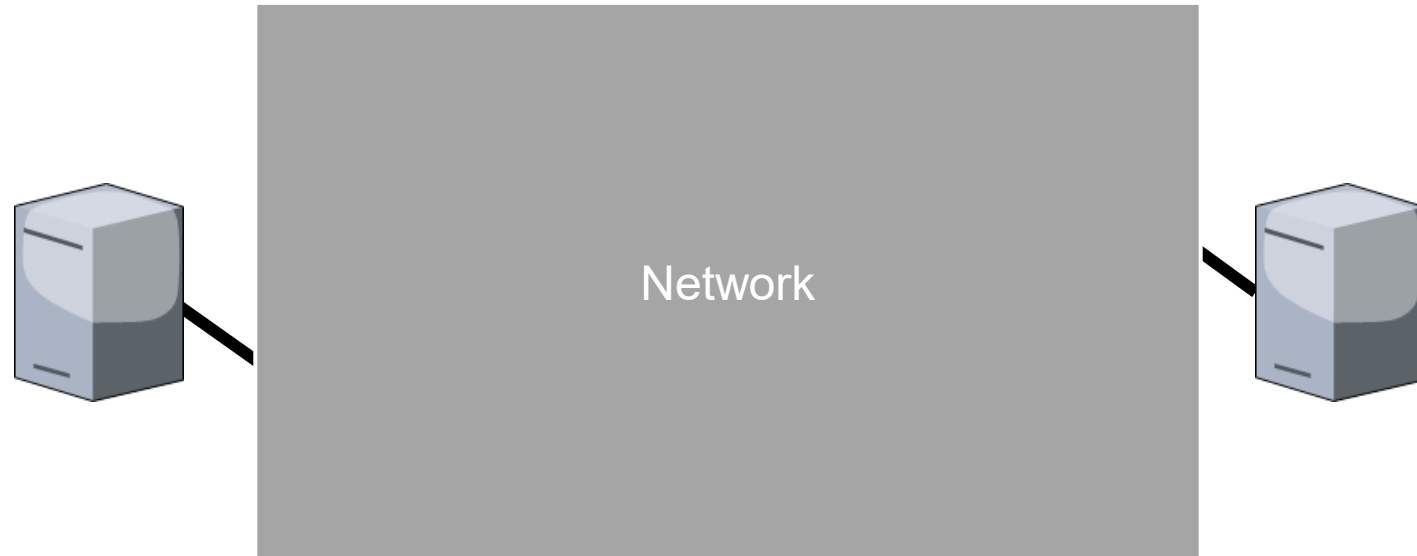
## Provided services

Runs only on the host  
and destination

Provides a **reliable** data stream over an **unreliable** network.

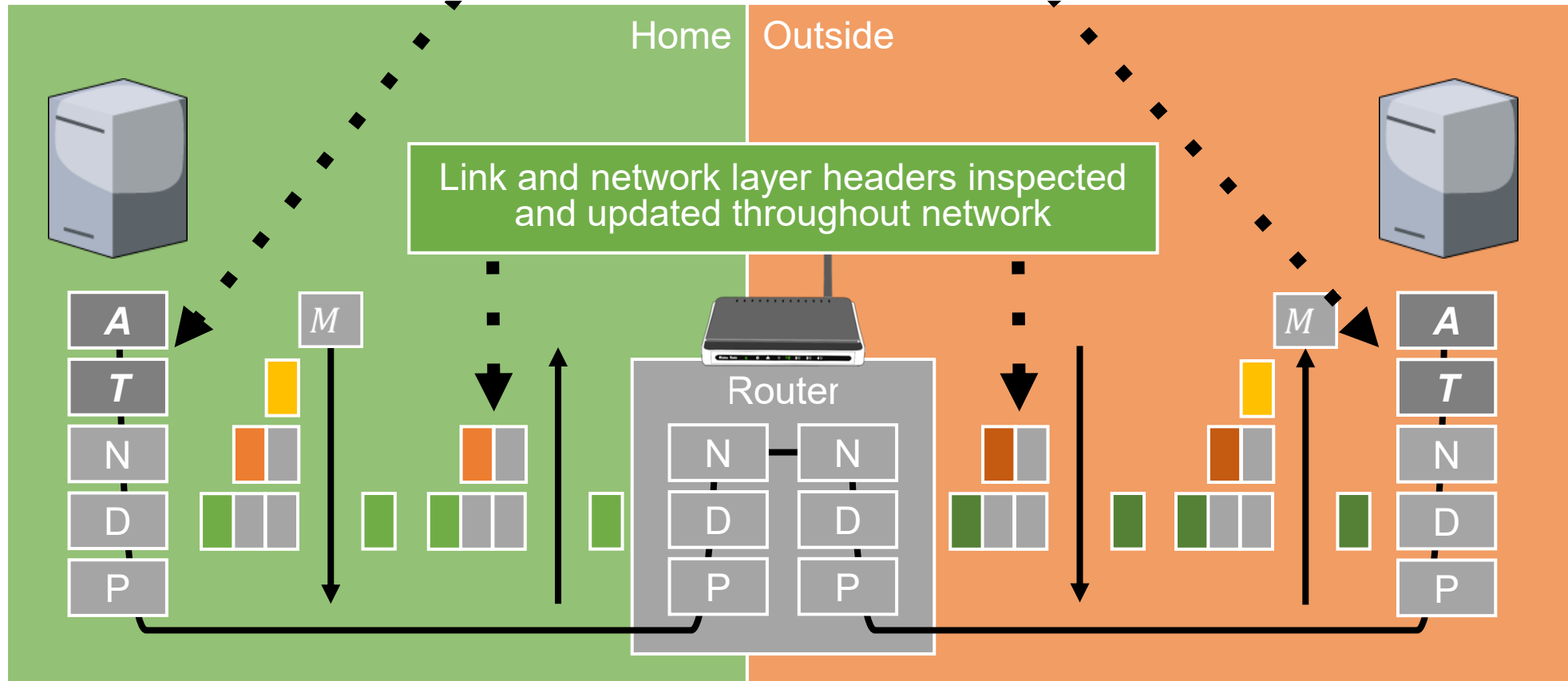
Provides communication between **applications**.

End-to-End Argument



# Transport layer only present at source and destination

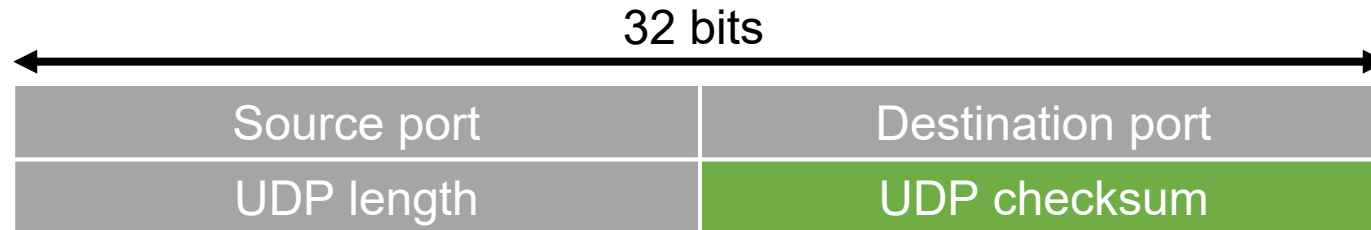
Transport layer and up used only at endpoints



# User Datagram Protocol (UDP)

RFC 768

**Very thin** layer on top of IP. Header provides **ports** needed to connect to remote applications.



The UDP header



Includes fields from the IP header!

UDP does **not** do:

1. Flow control
2. Congestion control
3. Retransmissions

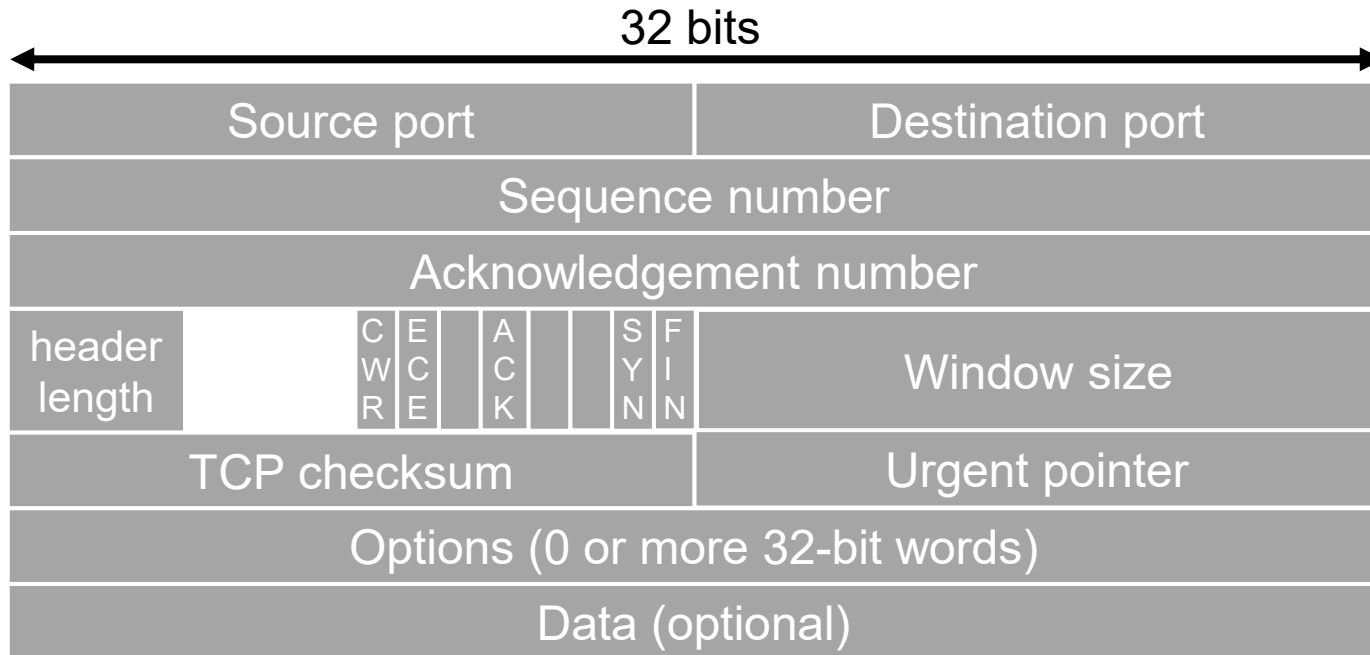
Q: Can you name a service that works well with UDP?

# Transmission Control Protocol (TCP)

One of the most important protocols on the Internet

Provides a **reliable end-to-end byte stream** over an unreliable network

Header: 20-60 bytes (Options: 0-40 bytes)



The TCP header

Q: How to calculate total TCP segment size?

Q: When are packets with 0 data bytes used?

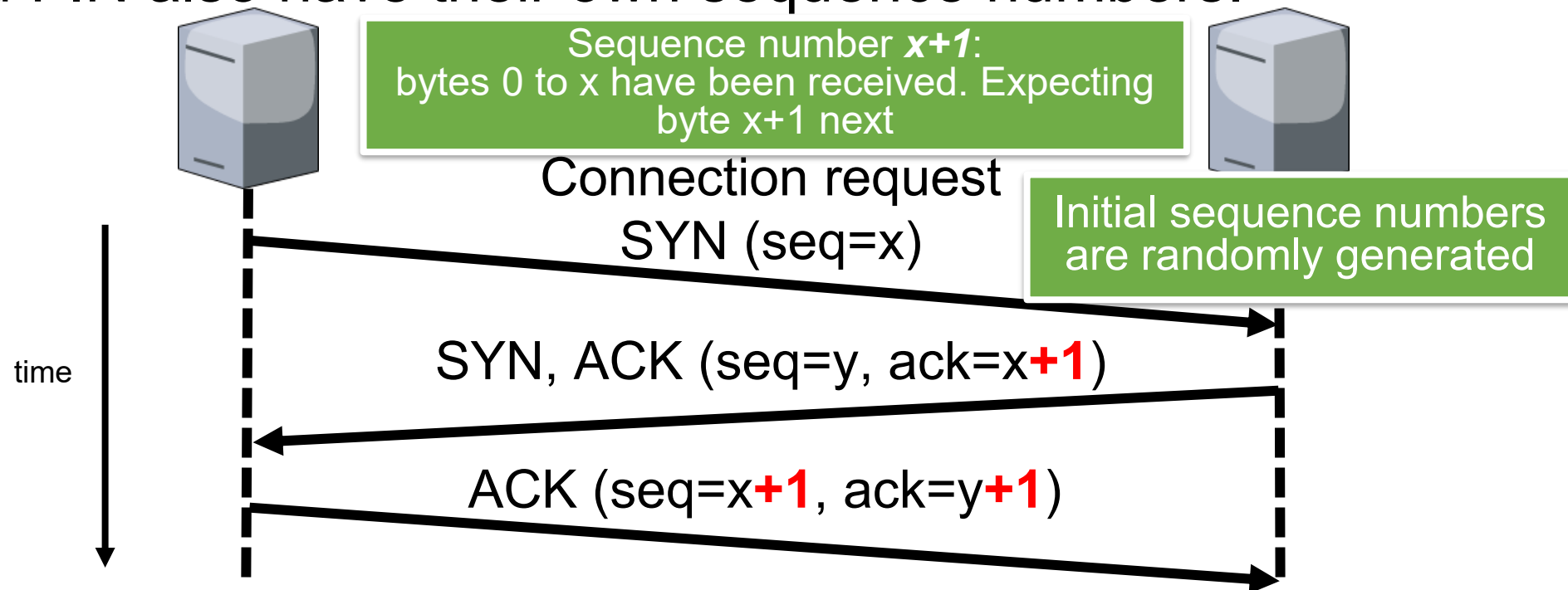
# TCP connection establishment

## Three-way handshake

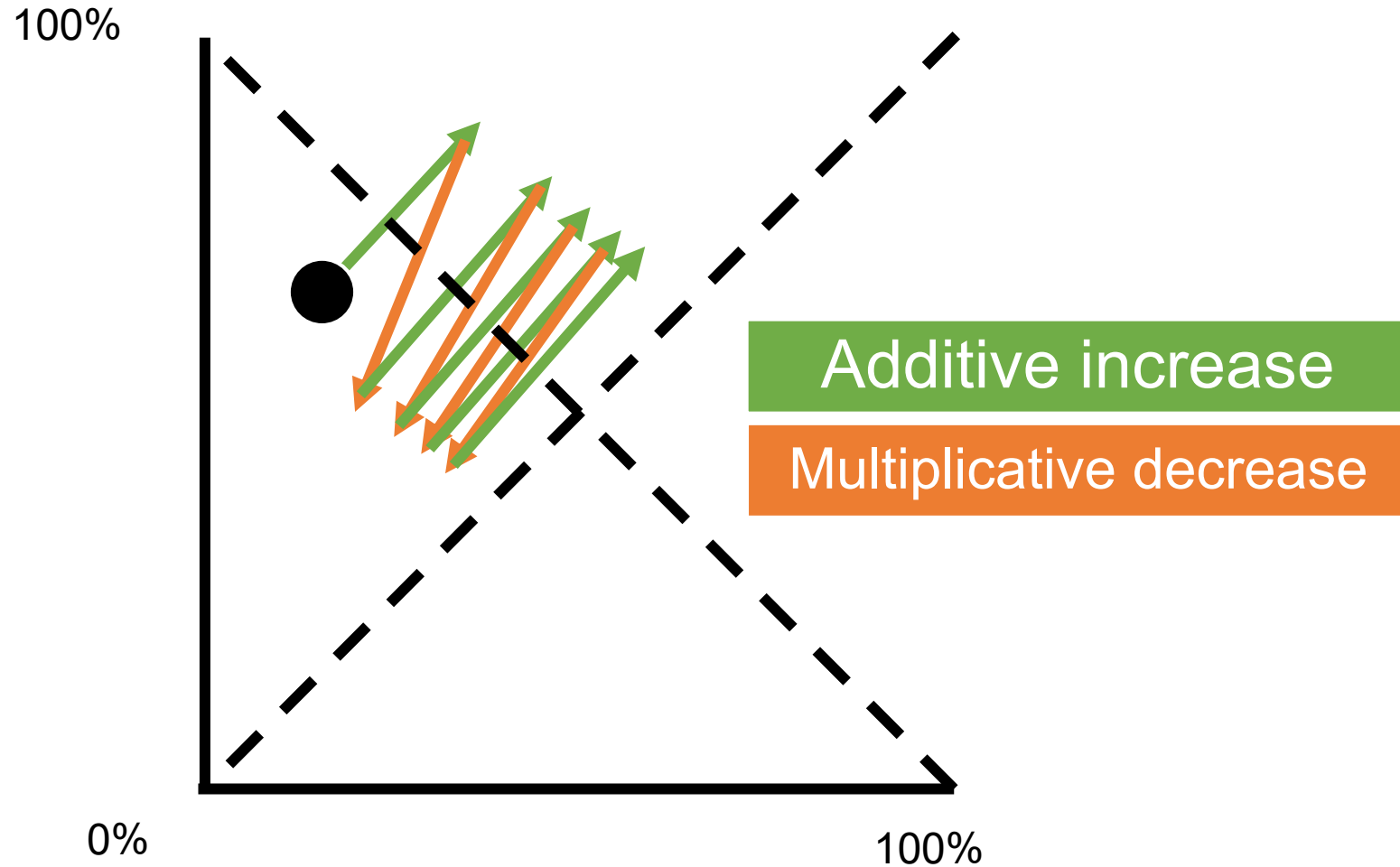
Uses timestamp option to improve performance on high-bandwidth networks

Every *data byte* has its own sequence number.\*

\*SYN and FIN also have their own sequence numbers.

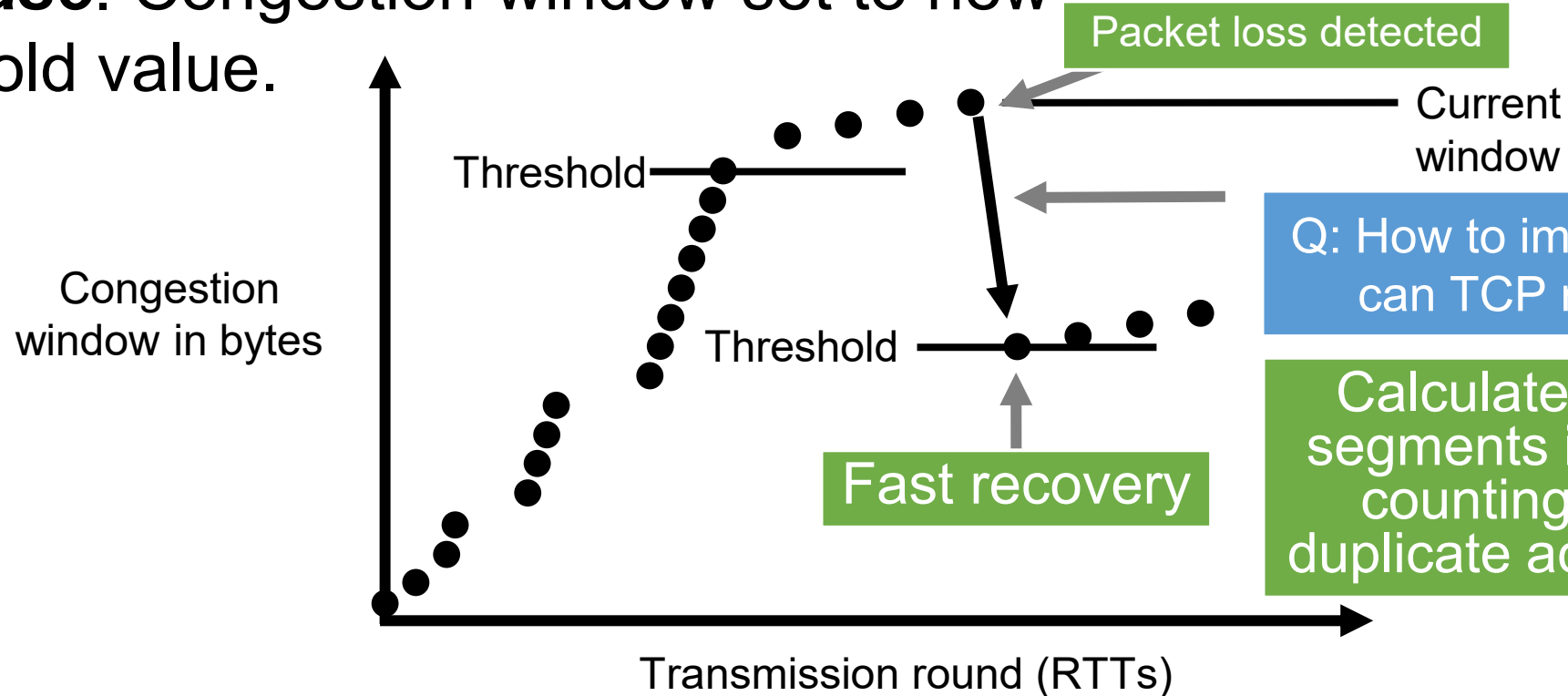


# Additive increase Multiplicative decrease



# TCP Reno (= TCP Tahoe + fast *recovery*)

Threshold reduced using *multiplicative decrease*. Congestion window set to new threshold value.



Q: How to implement this? When can TCP resume sending?

Calculates the number of segments in the network by counting the number of duplicate acknowledgements

# Recap: Link Utilization

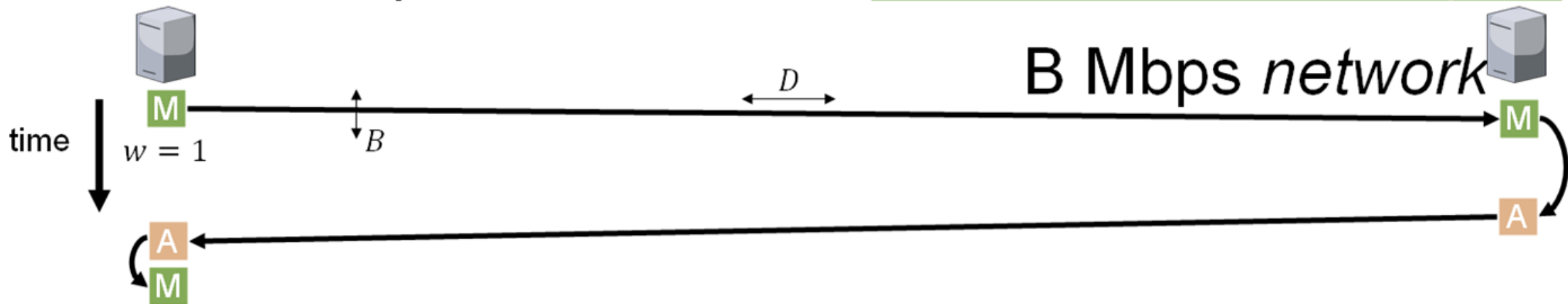
- Frame size (in bits/bytes):  $f$
- Window size (in frames):  $w$
- Bandwidth (max. data rate of physical channel):  $B_p$
- Bandwidth (frames per second):  $B_f$
- Propagation delay (in seconds):  $D$

It takes  $\frac{f}{B_p}$  seconds to send frame,  $\frac{B_p}{f} = B_f$

It takes  $D$  s for the frame to arrive at the receiver, takes  $D$  s for the (0-bit) acknowledgment to come back at the sender

1 frame per  $\frac{f}{B_p} + 2 \times D$  seconds

$$\text{Link utilization} = \frac{w}{1 + 2B_f D}$$



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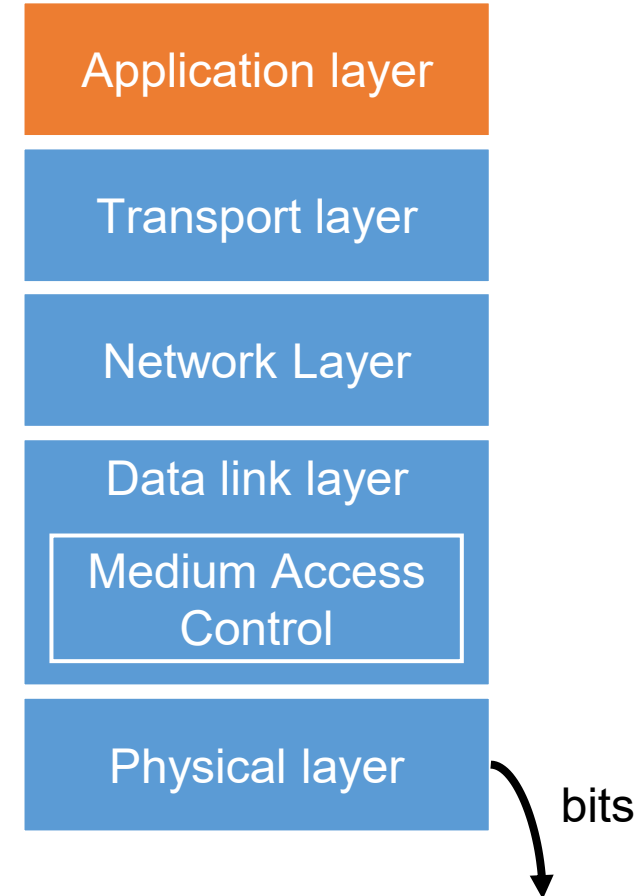
3. **Game demos**

4. **Wrap-up**

# Responsibilities of the Application Layer

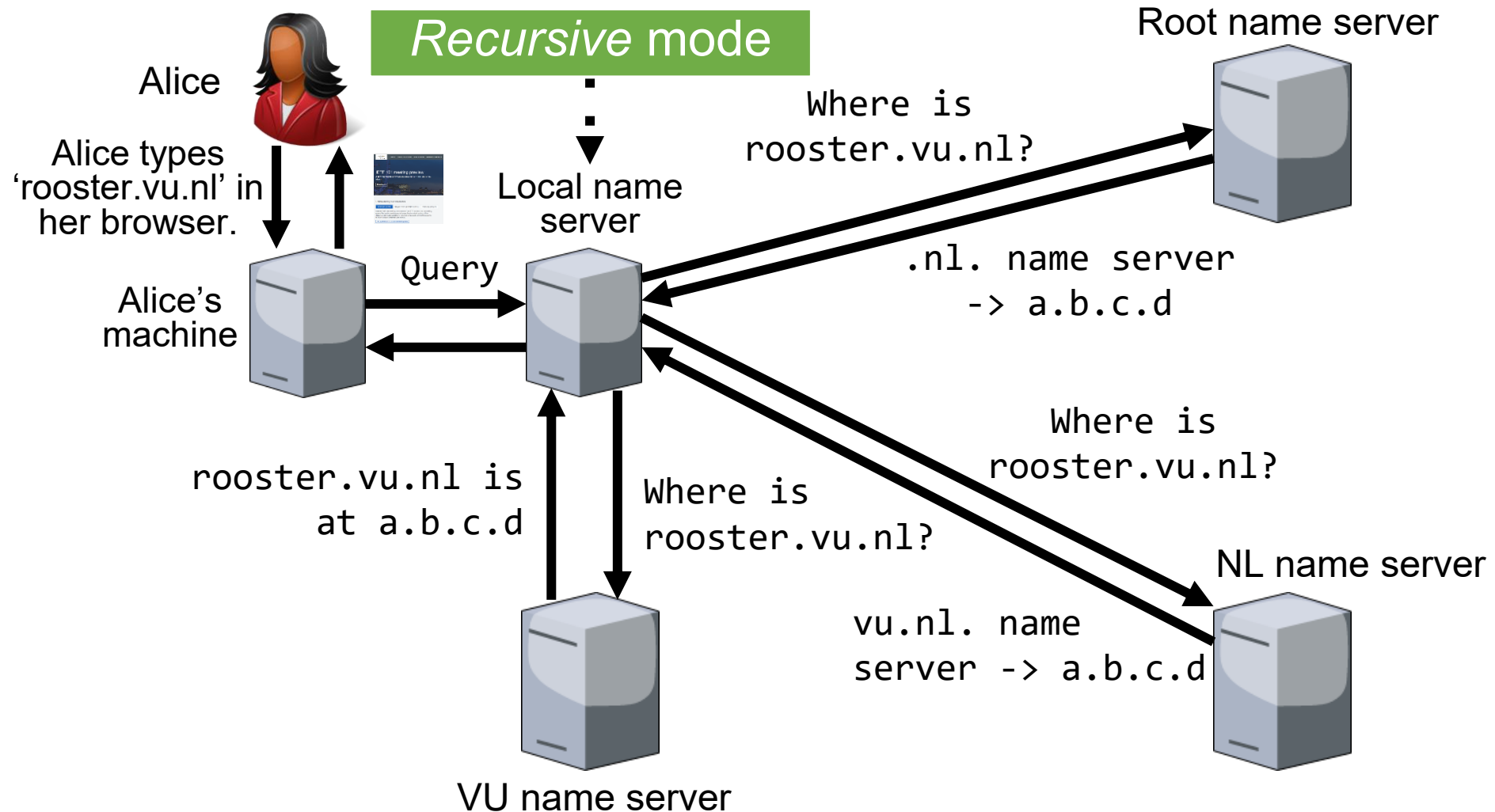
Provide applications that are (indirectly) useful to people

1. DNS
2. Mail and messaging
3. The Web
4. Multimedia applications
5. *Your future app here?*



Other name servers are in *iterative* mode

# Recursive and iterative DNS queries



# Base64 encoding

Q: How large is the overhead of base64 encoding?

Used to convert binary data to and from ASCII.

Alphabet: [A-Za-z0-9+/]

6 bits are translated into 1 character.

Q: What is the base64 encoding of 0xC0FFEE?

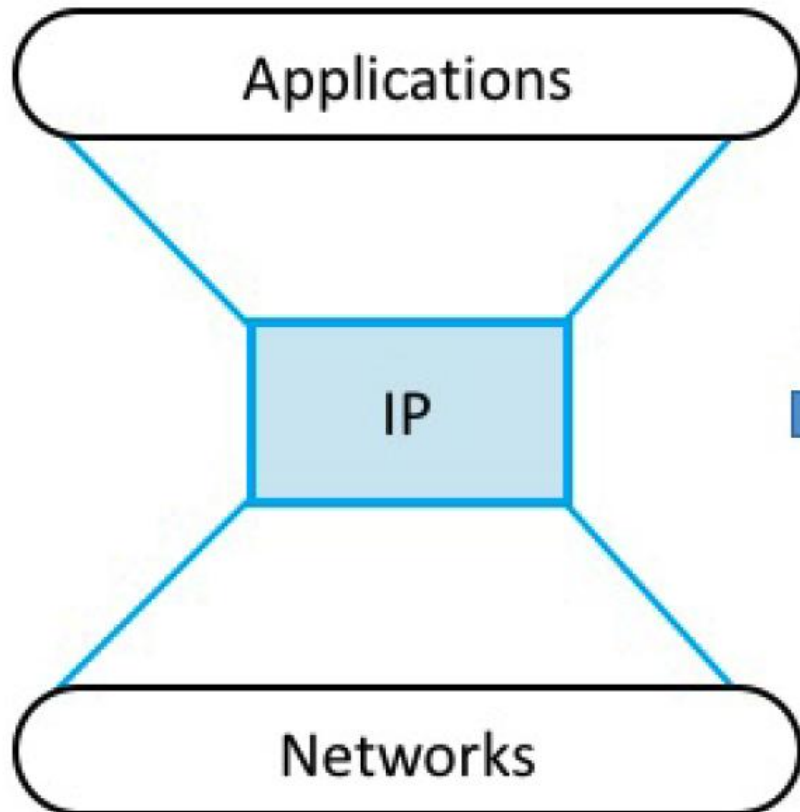
	C	a	t	
ASCII	<b>01000011</b>	01100001	<b>01110100</b>	
<hr/>				
Base64	<b>010000</b>	110110	<b>000101</b>	110100
	Q	2	F	0



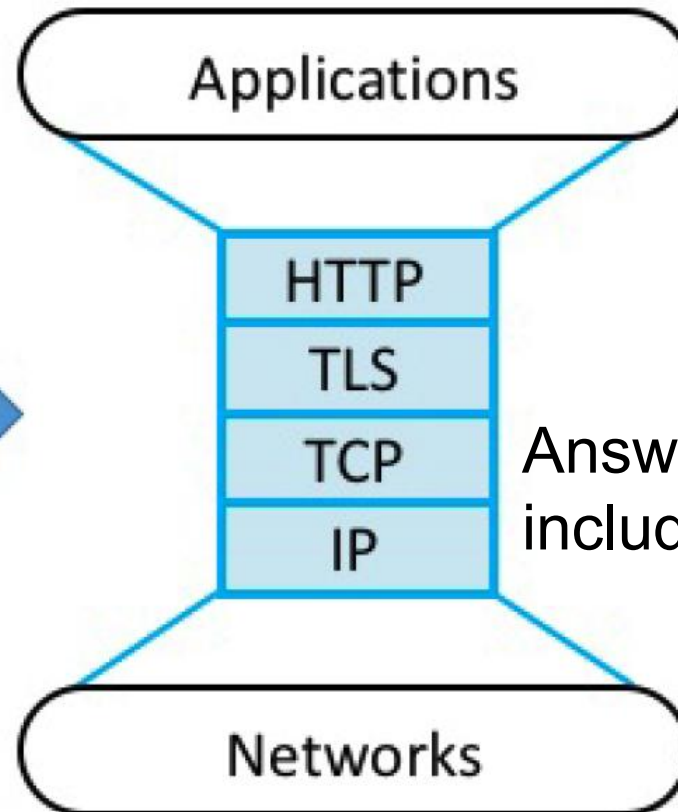
# HTTP is the new “narrow waist”

E.g., REST

Method	Description
GET	Read a Web page
HEAD	Read a Web page’s header
POST	Append to a Web page
PUT	Store a Web page



Evolution of the  
Narrow Waist



Answers include:

Q: Advantages over using TCP directly?

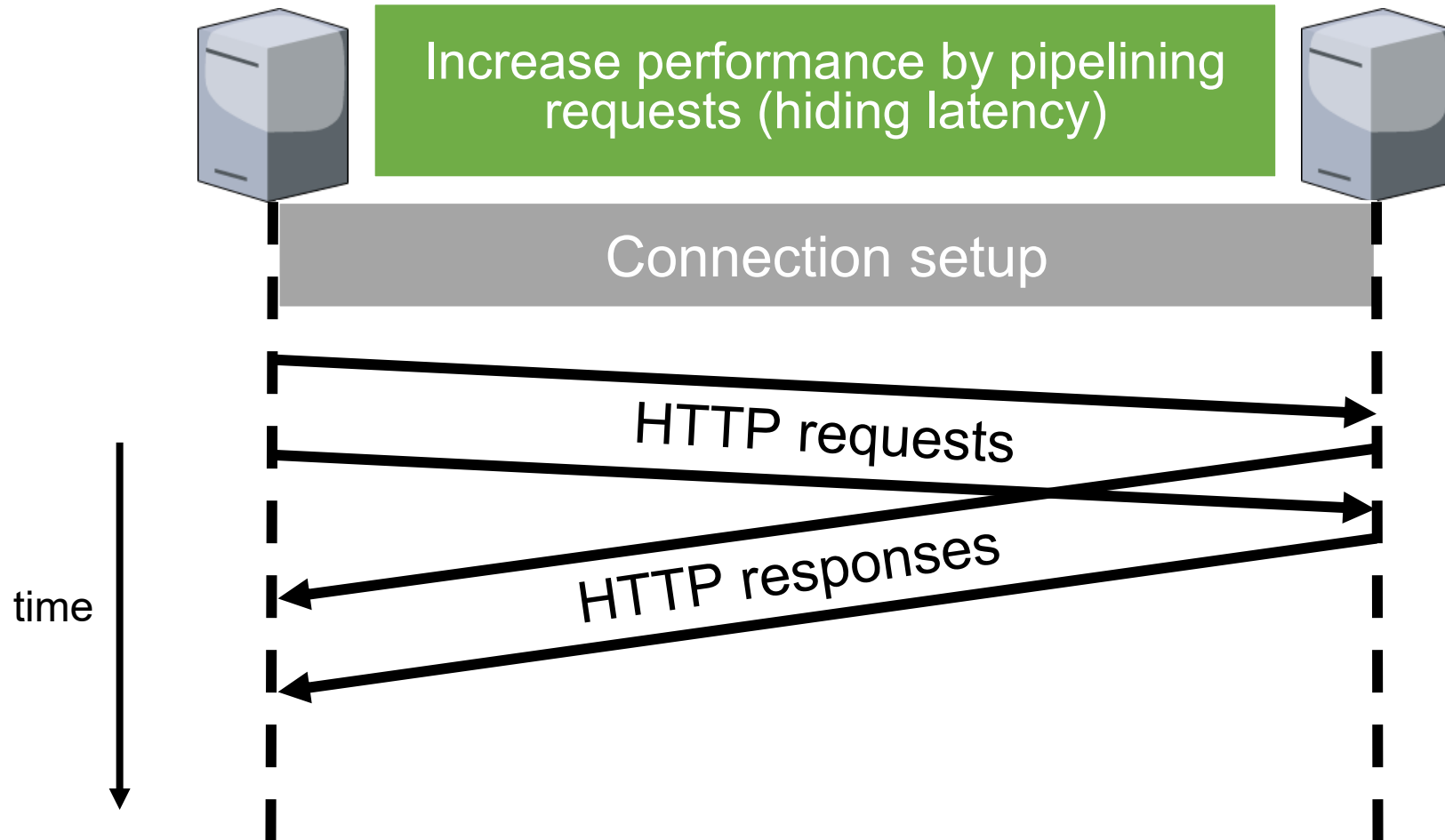
Provides set of methods

Provides security

Provides *naming*

# HTTP

## Pipelined requests

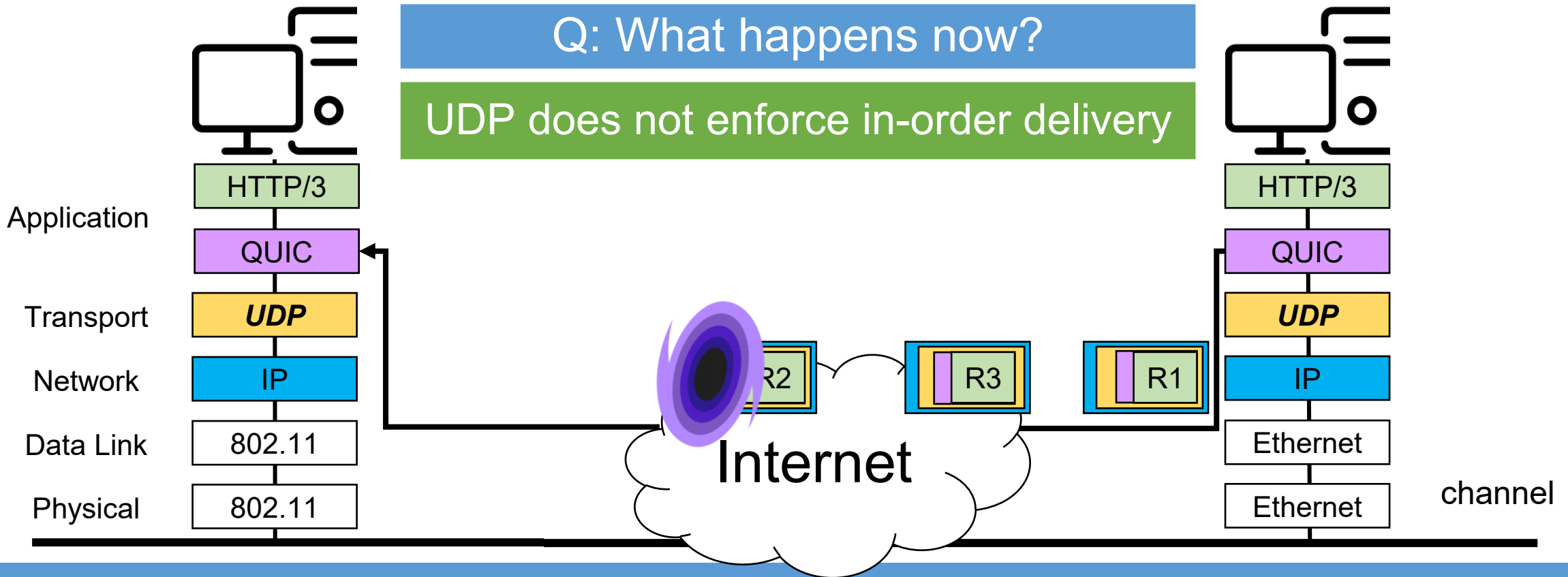


# HTTP/3 (HTTP + **QUIC**)

Each HTTP request can use a separate stream; within a stream, data is delivered in order; across streams no such guarantee is made

HTTP/3 uses the **QUIC** protocol

QUIC performs multiplexing, uses UDP



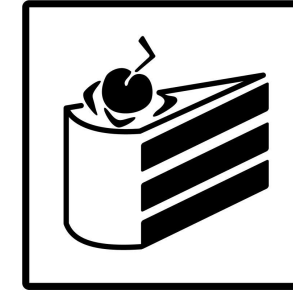
# Milestone reached!

Creating large-scale distributed systems is difficult!

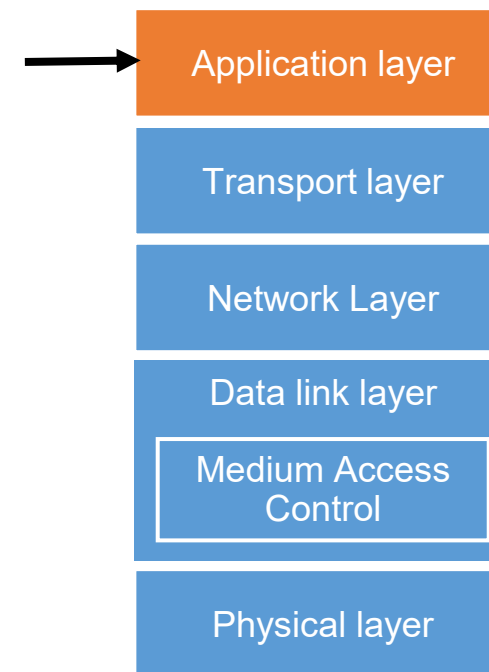
We can now **start** building applications and systems that communicate over a network!

Advanced courses unlocked:

1. Advanced Systems Programming
2. Advanced Computer Networks
3. Distributed systems  
(also requires Computer Organization and Operating Systems)



You are  
here



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**We hope you enjoyed the  
course and learned a lot!**